UNINFORMED SEARCH

Today

- Uninformed search
 - Formulating the search problem
 - State-space search
 - Analyze complexity of search

State-space search

- Search for a solution, i.e. a sequence of actions that leads from the initial state to the goal state
- Uninformed search algorithms
 - Uses no information beyond problem
 - Assumes a discrete environment
 - Offline exploration

Step One: Formulate the search problem

A well-defined search problem includes:

- states
- □ initial state

Induce the state space graph

- actions
- □ successor function
- □ goal test
- path cost (reflects performance measure)

Step One: Path to Bucharest



□ What does the state space graph look like?

Step One: 8-puzzle

5

6

7

4

1

3

Start State

8

2

- □ states?
- □ initial state?
- □ actions?
- □ successor function?
- goal test?
- □ path cost?
- □ What does the state space look like?

1	2	3
8		4
7	6	5

Goal State





Step Two: Search

- □ If actions are reversible, redundancy in search tree
- TreeSearch
 - Does not keep track of explored nodes
 - Infinite search tree
- GraphSearch
 - Keeps track of explored nodes
 - Search tree limited to size of state space

Step Two: GraphSearch

Search Strategies

A search strategy specifies the order in which nodes are selected from the frontier to be expanded

Breadth-first search (BFS)

- Expand shallowest unexpanded node
- □ Implementation: frontier is a FIFO queue



Breadth-first search (BFS)

Analyzing search algorithms

4 criteria for analyzing algorithms on board:

Notation

- □ b branching factor, i.e. max number of successors of any node
- □ d depth of the shallowest goal node
- m maximum length of any path in state space



Analyzing BFS

Depth-first search (DFS)

- Expand deepest unexpanded node
- Implementation: frontier is a LIFO queue (stack)



Analyzing DFS

Uniform-cost search

- Expand node with lowest path cost
- □ Implementation:
 - frontier is a priority queue ordered by path cost



Uniform-cost search

Analyzing Uniform-cost search

- $\hfill\square$ Let C* be the cost of the optimal solution and $\ensuremath{\mathcal{E}}$ be the minimum action cost
- □ Time: O(b^{C*/ ε + 1})
- □ Space: O($b^{C^*/\varepsilon + 1}$)
- □ Complete = YES if action cost exceeds epsilon
- Optimal = YES

Depth limited DFS

- DFS, but with a depth limit **L** specified
 - \blacksquare Nodes at depth ${\bm L}$ are treated as if they have no successors
 - We only search down to depth L
- □ Time?
 - □ O(b^L)
- Space?
 - O(bL)
- Complete?
 - No, if solution is longer than L
- Optimal
 - No, for same reasons DFS isn't

Iterative deepening search (IDS)

for L=0, 1, 2, ...

run depth-limited DFS with depth limit L if solution found return result

- Blends the benefits of BFS and DFS
 - searches in a similar order to BFS
 - but has the memory requirements of DFS
- Will find the solution when L is the depth of the shallowest goal

Iterative deepening search (IDS)



Time Complexity for IDS

Analysis of IDS

- 🗆 Time
 - □ O(b^d)
- Space
 - □ O(bd)
- □ Complete?
 - Yes
- Optimal?
 - Yes