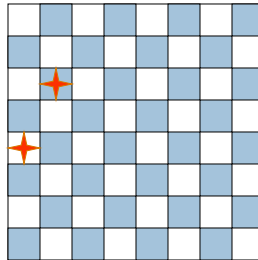


# CONSTRAINT SATISFACTION

## Today

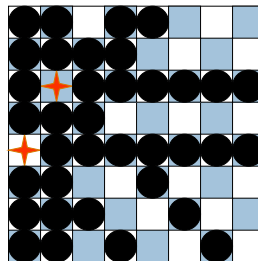
- Reading
  - AIMA Chapter 6
  
- Goals
  - Constraint satisfaction problems (CSPs)
  - Types of CSPs
  - Inference
  - Search + Inference

## 8-queens problem



How would you go about deciding where to put a third queen on the board in column 3?

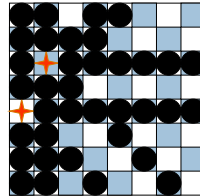
## 8-queens problem



How would you go about deciding where to put a third queen on the board in column 3?

## 8-queens problem

- This problem includes a set of **constraints**
- As a result, we need more than just a successor function and goal test
- We need a way to **propagate the constraints** imposed by one queen to the others and a way to detect **early failure**
  - ▣ Explicitly represent constraints
  - ▣ Algorithm to manipulate constraints

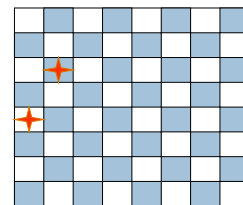


## Constraint satisfaction problems

- **Set of variables**  $\{X_1, X_2, \dots, X_n\}$
- Each variable  $X_i$  has a **domain**  $D_i$  of possible values
- **Set of constraints**  $\{C_1, C_2, \dots, C_p\}$ 
  - ▣ Each constraint  $C_k$  involves a subset of variables and specifies the allowable combinations of values to these variables
- A **state** is an assignment of values to some or all of the variables
  - ▣ If the assignment doesn't violate any constraints we say it is **consistent** or **legal**
- The **goal test** is checking for a consistent and complete assignment

## Example: 8-queens problems

- Variable?
- Domain?
- Constraints?

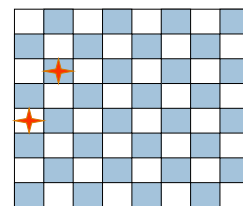


## Example: 8-queens problems

- **Variables:** one for each queen  $\{X_1, \dots, X_8\}$
- **Domain:** indicates row  $D = \{1, 2, \dots, 8\}$
- **Constraints:**

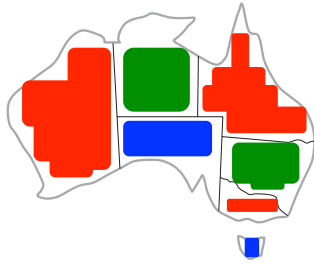
$$X_i = k \implies X_j \neq k \quad \forall i \neq j$$

$$X_i = k_i, X_j = k_j \implies |i - j| \neq |k_i - k_j|$$



## Example: Map coloring

- **Variables:** {WA, NT, SA, Q, NSW, V, T}
- **Domains:** {red, blue, green}
- **Constraints:** adjacent regions have different colors
  - ▣ Implicit:  $WA \neq NT, WA \neq SA, SA \neq NT, NT \neq Q, \dots$
  - ▣ Explicit:  $(WA, NT) \in \{(red, green), (red, blue), \dots\}$



## Example: Task scheduling

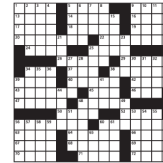
- **Variables:** {AxleF, AxleB, WheelRF, WheelLF, ..., Inspect}
- **Domains:** Time task starts  $D = [0, 1, 2, \dots, \infty)$
- **Constraints:**
  - ▣ Axle must be done before the wheel
    - $AxleF + 10 < WheelLF$
    - $AxleF + 10 < WheelRF$
  - ▣ The front axle and the back axle cannot be done at the same time
    - $(AxleF + 10 < AxleB)$  OR  $(AxleB + 10 < AxleF)$
  - ▣ Everything must be done within 30 minutes
    - Change domains to have upper bound 30 min.

## More examples

### More toy examples

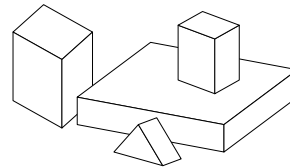
- sudoku, cryptarithmic

			2	8		7	
		3					8
	8			1			4
4					7		6
8		7	5	6		4	
5		7					1
9			8		6		
8				9			
2		5	4				



### Real-world applications

- Interpreting lines in 3D
- Assignment problems, e.g. who teaches what class?
- Timetable problems, e.g. which class offered when? where?
- Transportation scheduling
- Factory scheduling
- Circuit layout



## Types of CSPs - variables

### Discrete variables

#### Finite domains

- size  $d$  means  $O(d^n)$  possible assignments to explore

#### Infinite domains

- Linear constraints (e.g.  $T_1 + d_1 \leq T_2$ ) are solvable
- Non-linear constraints undecidable

### Continuous variables

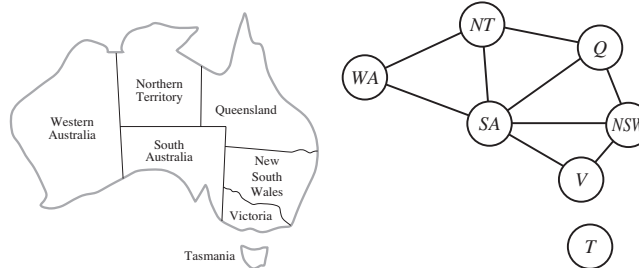
- linear programming problems with linear equality or inequality constraints solvable in polynomial time

## Types of CSPs - constraints

- **Unary constraints** involve a single variable
  - e.g. SA  $\neq$  green
- **Binary constraints** involve pairs of variables
  - SA  $\neq$  NSW
  - A binary CSP can be illustrated using a constraint graph
- **Higher-order constraints**
  - e.g. A, B, and C cannot be in the same grouping
  - e.g. AllDiff (all variables must be assigned different values)
- **Preference constraints**
  - costs on individual variable assignments
  - constraint optimization problem

## Constraint Graph

- Useful for **binary constraint CSPs** where each constraint relates (at most) two variables
- Nodes correspond to variables
- Edges (**arcs**) link two variables that participate in a constraint
- Use graph to speed up search



## Solving CSPs: Constraint Propagation

- Use the constraints to reduce the number of legal values for a variable
- Possible to find a solution without searching
  - Node consistency
    - A node is **node-consistent** if all values in its domain satisfy the unary constraints
  - Arc consistency
    - A node  $X_i$  is **arc-consistent** w.r.t. node  $X_j$  if for every value in  $D_i$  there exists a value in  $D_j$  that satisfies the binary constraint
    - Algorithm AC-3
  - Other types of consistency (path consistency, k-consistency, global constraints)

## AC-3 algorithm for Arc consistency

**function** AC-3(*csp*) **returns** false if inconsistency found, true otherwise

```

queue ← all arcs in csp
while queue not empty
  (Xi, Xj) ← REMOVE-FIRST(queue)
  if REMOVE-INCONSISTENT-VALUES(Xi, Xj)
    if size Di == 0 return false
    for each arc (Xk, Xi)
      add (Xk, Xi) to queue
return true
  
```

**function** REMOVE-INCONSISTENT-VALUES(X<sub>i</sub>, X<sub>j</sub>)

```

revised ← false
for each x in Di
  if ∄ y in Dj s.t. (x,y) satisfies constraints
    delete x from Di
  revised ← true
return revised
  
```



## AC-3 algorithm for Arc consistency

```

function AC-3(csp) returns false if inconsistency found, true otherwise
  queue ← all arcs in csp
  while queue not empty
    (Xi, Xj) ← REMOVE-FIRST(queue)
    if REMOVE-INCONSISTENT-VALUES(Xi, Xj)
      if size Di == 0 return false
      for each arc (Xk, Xi)
        add (Xk, Xi) to queue
  return true
  
```

$c$  constraints (arcs)  
 $d$  domain size  
  
 Total  
 $O(cd^3)$

---

```

function REMOVE-INCONSISTENT-VALUES(Xi, Xj)
  revised ← false
  for each x in Di
    if  $\nexists$  y in Dj s.t. (x, y) satisfies constraints
      delete x from Di
    revised ← true
  return revised
  
```

$O(d^2)$