

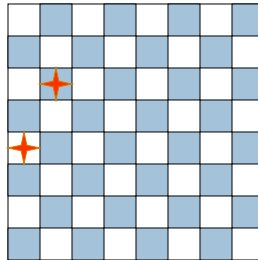
CONSTRAINT SATISFACTION

Today

- Reading
 - AIMA Chapter 6

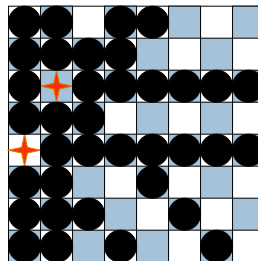
- Goals
 - Constraint satisfaction problems (CSPs)
 - Types of CSPs
 - Inference
 - Search + Inference

8-queens problem



How would you go about deciding where to put a third queen on the board in column 3?

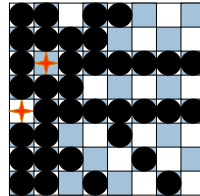
8-queens problem



How would you go about deciding where to put a third queen on the board in column 3?

8-queens problem

- This problem includes a set of **constraints**
- As a result, we need more than just a successor function and goal test
- We need a way to **propagate the constraints** imposed by one queen to the others and a way to detect **early failure**
 - ▣ Explicitly represent constraints
 - ▣ Algorithm to manipulate constraints

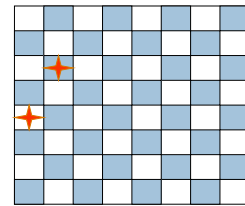


Constraint satisfaction problems

- **Set of variables** $\{X_1, X_2, \dots, X_n\}$
- Each variable X_i has a **domain** D_i of possible values
- **Set of constraints** $\{C_1, C_2, \dots, C_p\}$
 - ▣ Each constraint C_k involves a subset of variables and specifies the allowable combinations of values to these variables
- A **state** is an assignment of values to some or all of the variables
 - ▣ If the assignment doesn't violate any constraints we say it is **consistent** or **legal**
- The **goal test** is checking for a consistent and complete assignment

Example: 8-queens problems

- Variable?
- Domain?
- Constraints?

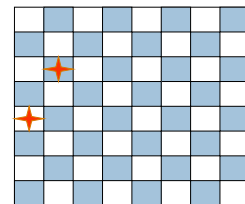


Example: 8-queens problems

- **Variables:** one for each queen $\{X_1, \dots, X_8\}$
- **Domain:** indicates row $D = \{1, 2, \dots, 8\}$
- **Constraints:**

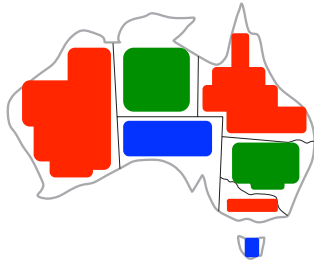
$$X_i = k \implies X_j \neq k \quad \forall i \neq j$$

$$X_i = k_i, X_j = k_j \implies |i - j| \neq |k_i - k_j|$$



Example: Map coloring

- **Variables:** {WA, NT, SA, Q, NSW, V, T}
- **Domains:** {red, blue, green}
- **Constraints:** adjacent regions have different colors
 - ▣ Implicit: $WA \neq NT, WA \neq SA, SA \neq NT, NT \neq Q, \dots$
 - ▣ Explicit: $(WA, NT) \in \{(red, green), (red, blue), \dots\}$



Example: Task scheduling

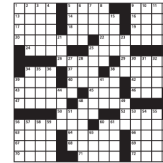
- **Variables:** {AxleF, AxleB, WheelRF, WheelLF, ..., Inspect}
- **Domains:** Time task starts $D = [0, 1, 2, \dots, \infty)$
- **Constraints:**
 - ▣ Axle must be done before the wheel
 - $AxleF + 10 < WheelLF$
 - $AxleF + 10 < WheelRF$
 - ▣ The front axle and the back axle cannot be done at the same time
 - $(AxleF + 10 < AxleB)$ OR $(AxleB + 10 < AxleF)$
 - ▣ Everything must be done within 30 minutes
 - Change domains to have upper bound 30 min.

More examples

More toy examples

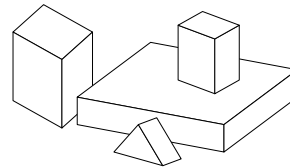
- ▣ sudoku, cryptarithmic

			2	8		7	
		3					8
	8			1			4
4					7		6
8		7	5	6		4	
5		7					1
9			8		6		
8				9			
2		5	4				



Real-world applications

- ▣ Interpreting lines in 3D
- ▣ Assignment problems, e.g. who teaches what class?
- ▣ Timetable problems, e.g. which class offered when? where?
- ▣ Transportation scheduling
- ▣ Factory scheduling
- ▣ Circuit layout



Types of CSPs - variables

Discrete variables

Finite domains

- ▣ size d means $O(d^n)$ possible assignments to explore

Infinite domains

- ▣ Linear constraints (e.g. $T_1 + d_1 \leq T_2$) are solvable
- ▣ Non-linear constraints undecidable

Continuous variables

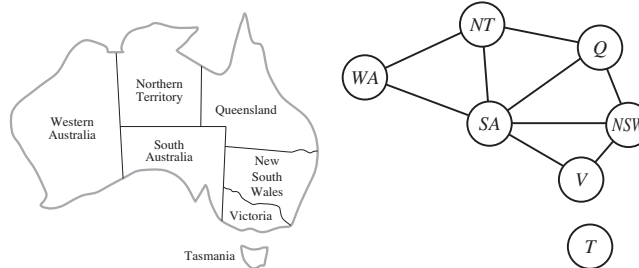
- ▣ linear programming problems with linear equality or inequality constraints solvable in polynomial time

Types of CSPs - constraints

- **Unary constraints** involve a single variable
 - e.g. SA \neq green
- **Binary constraints** involve pairs of variables
 - SA \neq NSW
 - A binary CSP can be illustrated using a constraint graph
- **Higher-order constraints**
 - e.g. A, B, and C cannot be in the same grouping
 - e.g. AllDiff (all variables must be assigned different values)
- **Preference constraints**
 - costs on individual variable assignments
 - constraint optimization problem

Constraint Graph

- Useful for **binary constraint CSPs** where each constraint relates (at most) two variables
- Nodes correspond to variables
- Edges (**arcs**) link two variables that participate in a constraint
- Use graph to speed up search



Inference: constraint propagation

- Use the constraints to reduce the number of legal values for a variable
- Possible to find a solution without searching
 - Node consistency
 - A node is **node-consistent** if all values in its domain satisfy the unary constraints
 - Arc consistency
 - A node X_i is **arc-consistent** w.r.t. node X_j if for every value in D_i there exists a value in D_j that satisfies the binary constraint
 - Algorithm AC-3
 - Other types of consistency (path consistency, k-consistency, global constraints)

AC-3 algorithm for Arc consistency

function AC-3(*csp*) **returns** false if inconsistency found, true otherwise

```

queue ← all arcs in csp
while queue not empty
  ( $X_i, X_j$ ) ← REMOVE-FIRST(queue)
  if REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ )
    if size  $D_i$  == 0 return false
    for each arc ( $X_k, X_i$ )
      add ( $X_k, X_i$ ) to queue
return true
  
```

function REMOVE-INCONSISTENT-VALUES(X_i, X_j)

```

revised ← false
for each x in  $D_i$ 
  if  $\nexists$  y in  $D_j$  s.t. (x,y) satisfies constraints
    delete x from  $D_i$ 
  revised ← true
return revised
  
```


AC-3 algorithm for Arc consistency

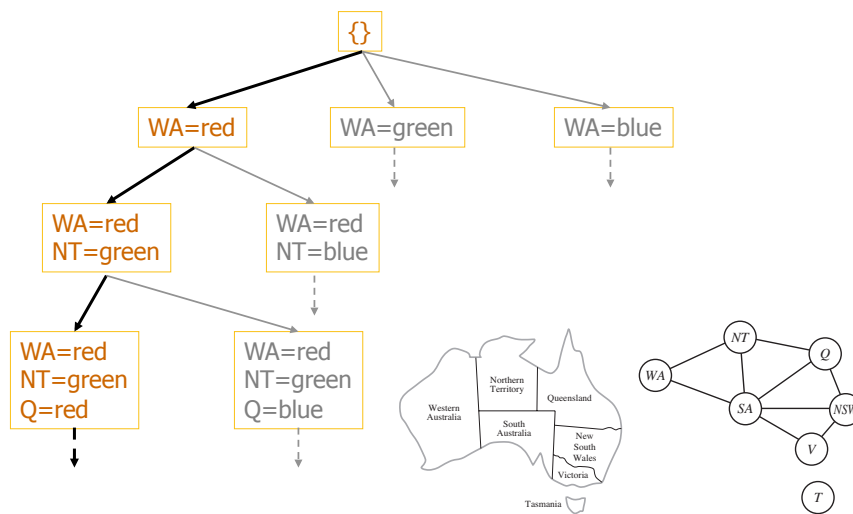
```

function AC-3(csp) returns false if inconsistency found, true otherwise
    queue ← all arcs in csp
    while queue not empty
        (Xi, Xj) ← REMOVE-FIRST(queue)
        if REMOVE-INCONSISTENT-VALUES(Xi, Xj)
            if size Di == 0 return false
            for each arc (Xk, Xi)
                add (Xk, Xi) to queue
    return true

function REMOVE-INCONSISTENT-VALUES(Xi, Xj)
    revised ← false
    for each x in Di
        if ∄ y in Dj s.t. (x, y) satisfies constraints
            delete x from Di
        revised ← true
    return revised
    
```

c constraints (arcs)
 d domain size
 $O(cd)$
 $O(d^2)$
 Total $O(cd^3)$

Search



Backtracking Search

```

function CSP-BACKTRACKING(assignment) returns a solution or failure
  if assignment complete return assignment
  X ← select unassigned variable
  D ← select an ordering for the domain of X
  for each value in D
    if value is consistent with assignment
      add (X = value) to assignment
      (ADD INFERENCE HERE)
      result ← CSP-BACKTRACKING(assignment)
      if result ≠ failure return result
    remove (X = value) from assignment
  return failure
  
```

Backtracking Search

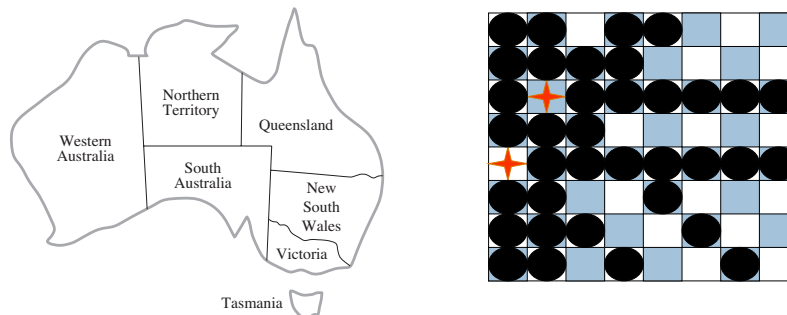
- Backtracking search is the basic uninformed algorithm for solving CSPs
- **Idea 1: One variable at a time**
 - Variable assignments are commutative so fix ordering
 - i.e. (WA = red, NT = green) is the same as (NT = green, WA = red)
- **Idea 2: Check constraints as we go**
 - Consider only values which do not conflict with previous assignments
 - May take some computation to check
 - “incremental goal test”
 - (Additional inference is optional, e.g. arc-consistency)
- Depth-first search with these 2 improvements is called *backtracking search*

Improving backtracking search

- Ordering
 - ▣ Which variable X should be assigned a value next?
 - ▣ In which order should its domain D be sorted?
- Incorporating inference
 - ▣ Can we detect inevitable failure early?
- Structure
 - ▣ Can we exploit the problem structure?

Ordering

- Which variable should we choose?



Ordering

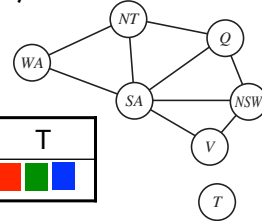
- Variable ordering
 - Minimum-remaining values heuristic - Choose the variable with the fewest “legal” moves remaining
 - Degree heuristic - Choose variable involved in the largest number of constraints with remaining unassigned variables
- Value ordering
 - Least-constraining value heuristic - Choose the value that rules out the fewest choices for the neighboring variables

Improving backtracking search

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Incorporating inference: forward checking

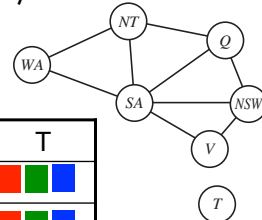
- After an assignment $X = x$, look at each unassigned variable Y connected to X by a constraint
 - ▣ Delete from Y 's domain any value inconsistent with $X = x$
 - ▣ Equivalent to ensuring all arcs of the form (Y, X) are arc consistent



WA	NT	Q	NSW	V	SA	T
Red, Green, Blue	Red, Green, Blue	Red, Green, Blue	Red, Green, Blue	Red, Green, Blue	Red, Green, Blue	Red, Green, Blue

Incorporating inference: forward checking

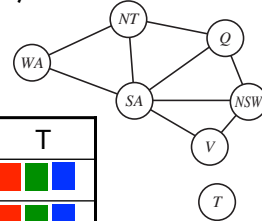
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Incorporating inference: forward checking

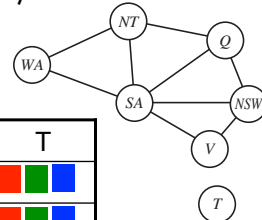
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Incorporating inference: forward checking

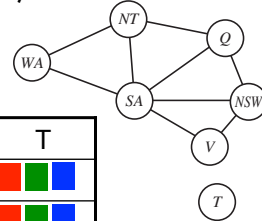
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Incorporating inference: forward checking

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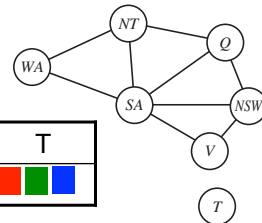


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Red	Blue	Green	Red, Blue	Red, Green, Blue	Blue	Red, Green, Blue
Red	Blue	Green	Red	Blue		Red, Green, Blue

Could have detected earlier that things were going wrong!

Incorporating inference: AC-3

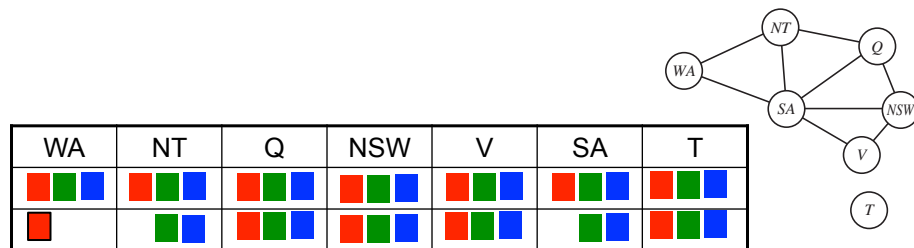
- Can also infer **path consistency** based on triples of variables and **k-consistency**



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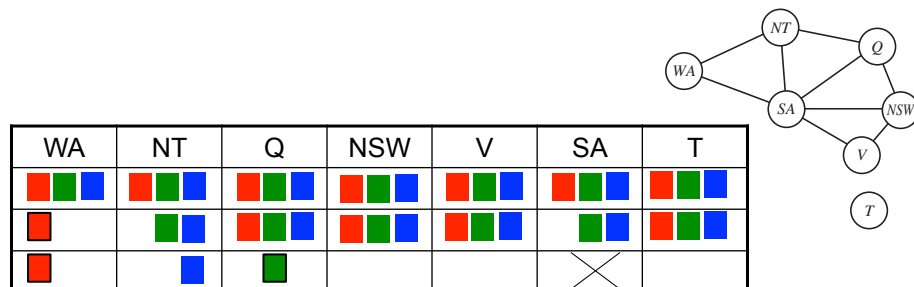
Incorporating inference: AC-3

- Can also infer **path consistency** based on triples of variables and **k-consistency**



Incorporating inference: AC-3

- Can also infer **path consistency** based on triples of variables and **k-consistency**

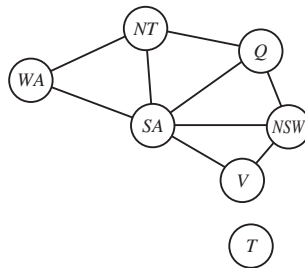


Improving backtracking search

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Structure of the constraint graph

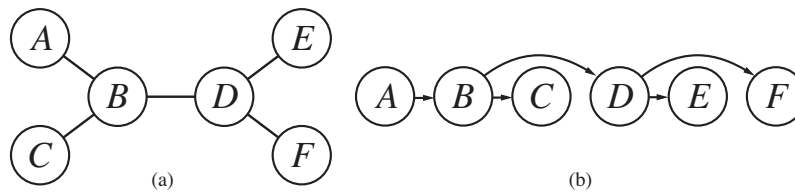
- Independent subproblems
 - ▣ Find connected components of the constraint graph
 - ▣ e.g. Tasmania and the mainland are independent
 - ▣ If we can split n variables into c subproblems of n/c variables each: $O(d^n) \longrightarrow O(d^c n/c)$



Tasmania is independent of the mainland!

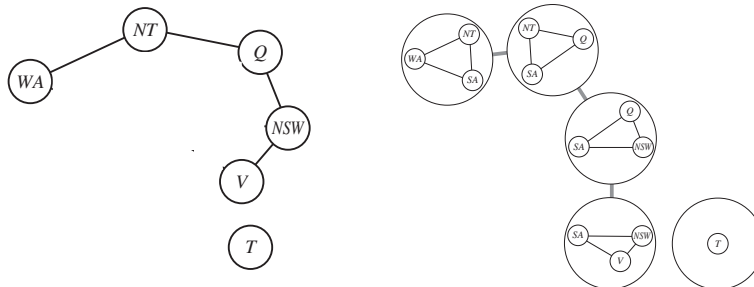
Structure of the constraint graph

- Tree structured constraint graphs
 - ▣ Can solve in linear time using AC-3



Structure of the constraint graph

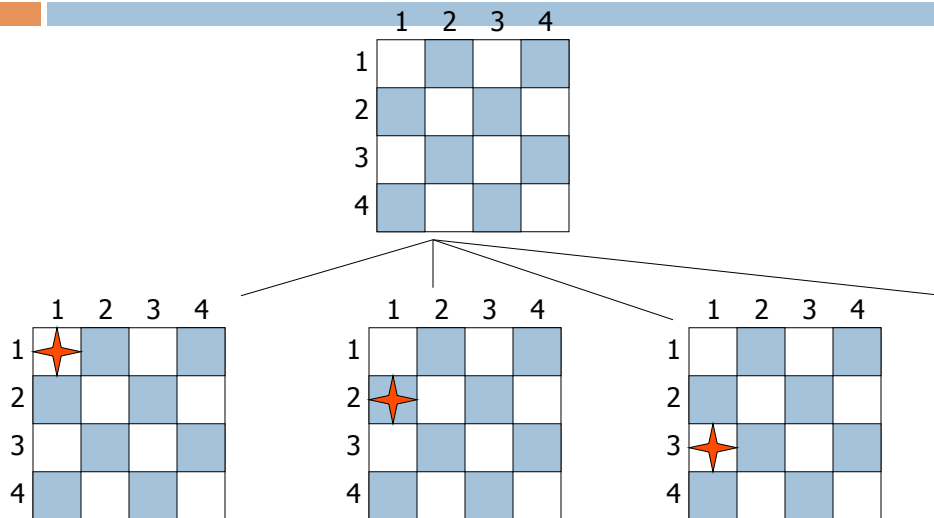
- Reduction to a tree structured graph
 - ▣ Cycle cutset – a subset of the variables whose removal creates a tree.
 - ▣ Tree decomposition – Divide graph into subproblems, solve independently merge the solutions



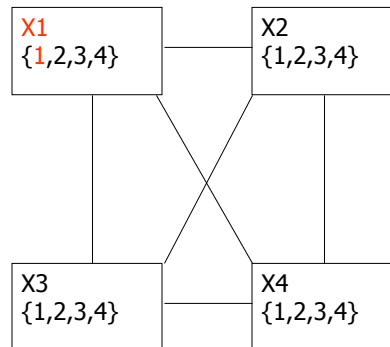
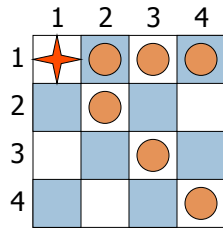
CSP Summary

- Constraint Satisfaction Problems (CSPs)
- Solving CSPs using inference
- Solving CSPs using search
 - ▣ Backtracking algorithm = DSF + constraints checking
 - ▣ General (not problem-specific) heuristics
- Improving Backtracking
 - ▣ Variable and value ordering
 - ▣ Inference
 - ▣ Structure

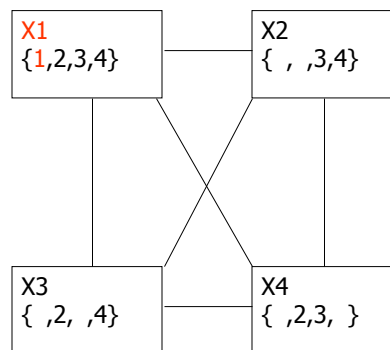
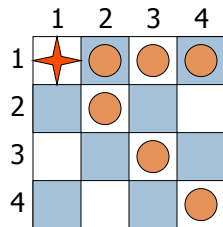
4-Queens Problem



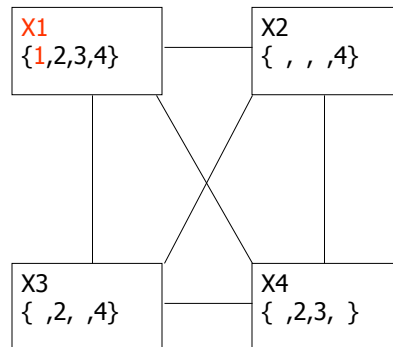
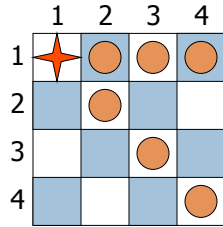
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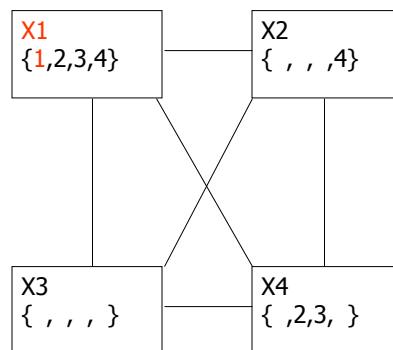
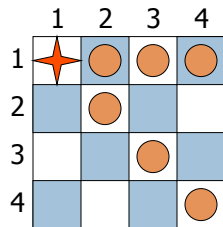
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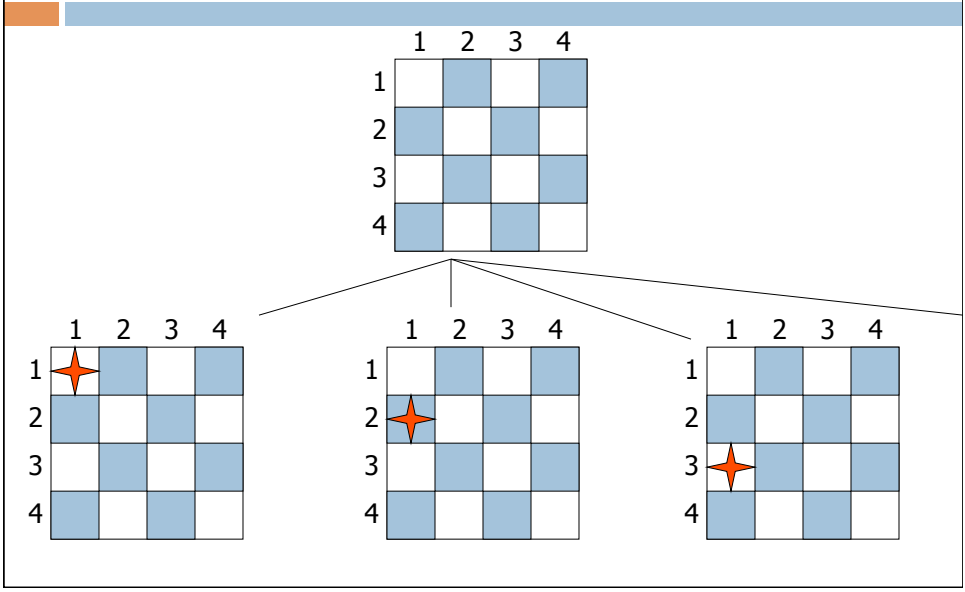
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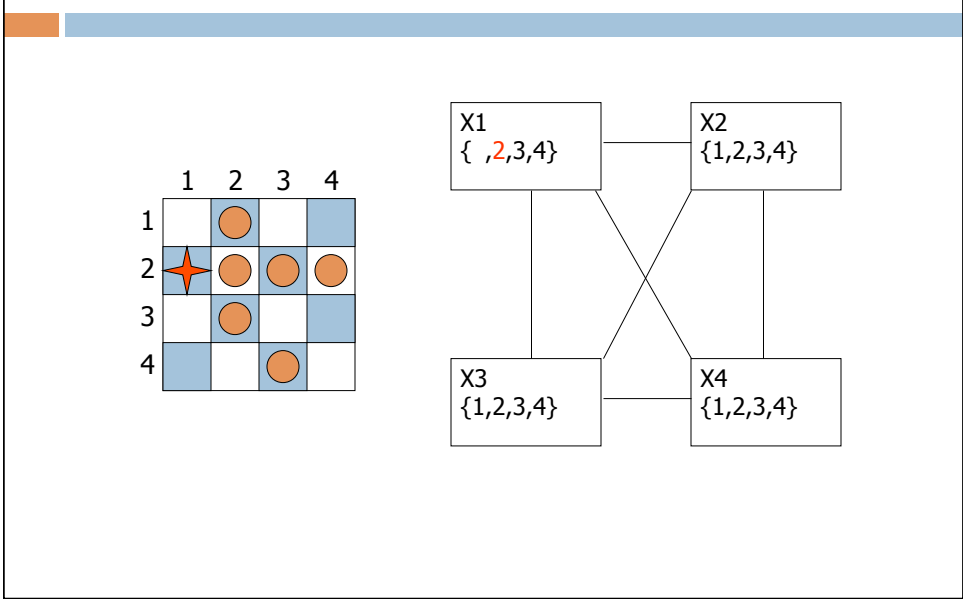
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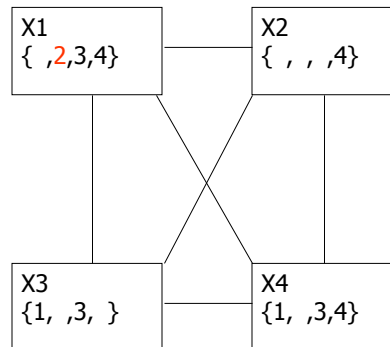
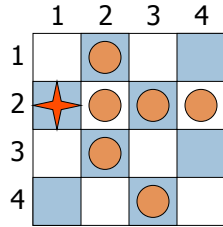
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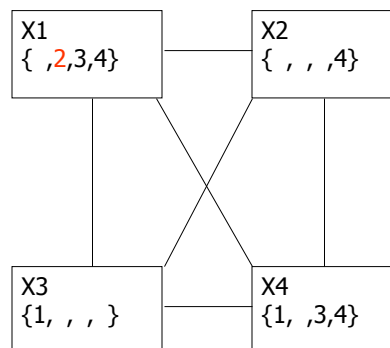
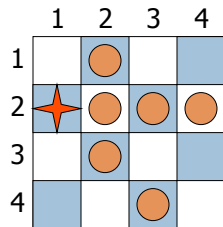
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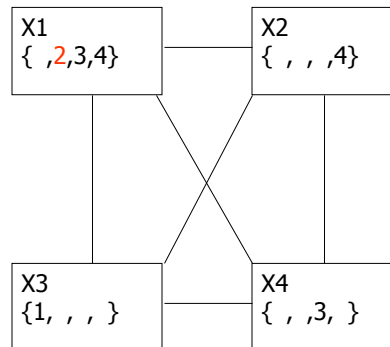
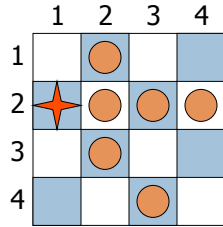
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