## ADVERSARIAL SEARCH



## Today

## Reading

- AIMA Chapter 5.1-5.5, 5.7,5.8

Goals

- alpha-beta pruning
$\square$ Finish real-time decisions
$\square$ Stochastic games


## Minimax: an optimal strategy


$\operatorname{MinimaX}(\mathrm{s})= \begin{cases}\operatorname{UTILITY}(s) & \text { if TERMINAL-TEST}(s) \\ \max _{a} \operatorname{MINIMAX}(\operatorname{ReSULT}(s, a)) & \text { if PLAYER }(s)=\operatorname{MAX} \\ \min _{a} \operatorname{MINIMAX}(\operatorname{RESULT}(s, a)) & \text { if } \operatorname{PLAYER}(s)=\operatorname{MIN}\end{cases}$

## Minimax: An Optimal Strategy

function Minimax-Decision(state) returns an action
$v \leftarrow \operatorname{MAX}-\operatorname{VALUE}($ state $)$
return the action in SUCCESSORS(state) with value $v$
function MAX-VALUE(state) returns a utility value
if Terminal-Test(state) then return Utlity(state)
$v \leftarrow-\infty$
for $a, s$ in Successors(state) do
$v \leftarrow \operatorname{Max}(v, \operatorname{Min}-\operatorname{VALUE}(s))$
return $v$
function Min-Value(state) returns a utility value
if Terminal-Test(state) then return Utlity(state)
$v \leftarrow \infty$
for $a, s$ in Successors(state) do
$v \leftarrow \operatorname{Min}(v, \operatorname{Max}-\operatorname{Value}(s))$
return $v$

## Minimax example



## Minimax example



Minimax example


## Minimax example



Minimax example


## Minimax example



Minimax example


## Complexity of Minimax

$\square$ Minimax performs DFS of search tree
$\square$ Time O(b $\left.{ }^{m}\right)$
$\square$ Tic-tac-toe: ~5 legal moves, 9 moves/game

- $5^{9}=1,953,125$ states
$\square$ Chess: $\sim 35$ legal moves, $\sim 100$ moves/game
- $35^{100}$ states to search

Common games produce enormous search trees

## Alpha-Beta pruning

alpha is the best scenario MAX has found so far
$\square$ MAX can always achieve a utility of alpha (and hopes for higher)
beta is the best scenario MIN has found so far
$\square$ MIN can always achieve a utility of beta (and hopes for lower)

$$
[\alpha, \beta]
$$

## Alpha-Beta Example

Do depth-first search until first leaf


## Alpha-Beta Example



## Alpha-Beta Example

MAX

Alpha-Beta Example


## Alpha-Beta Example



Alpha-Beta Example


## Alpha-Beta Example



Alpha-Beta Example

MAX



## Effectiveness of Alpha-Beta pruning

Highly-dependent on the order in which the states are examined

Try to examine those states that are likely to be best

In practice, the running time for alpha-beta pruning is $O\left(b^{m / 2}\right)$ as opposed to $O\left(b^{m}\right)$
$\square$ Effective branching factor is square root of $b$
$\square$ Can search twice as deep as minimax

## Real-time decision making

$\square$ Alpha-beta pruning still has to search down to the leaf nodes (for part of the search tree)

Standard approach (Shannon, 1950):
$\square$ apply a cutoff test (turn non-leaf nodes into leaves)
$\square$ replace utility function by an evaluation function that estimates "desirability" of position


Claude Shannon

## Real-time decision making

$\operatorname{MinimAX}(\mathrm{s})= \begin{cases}\operatorname{UTILITY}(s) & \text { if TERMINAL-TEST}(s) \\ \max _{a} \operatorname{MINIMAX}(\operatorname{RESULT}(s, a)) & \text { if } \operatorname{PLAYER}(s)=\operatorname{MAX} \\ \min _{a} \operatorname{MINIMAX}(\operatorname{RESULT}(s, a)) & \text { if } \operatorname{PLAYER}(s)=\operatorname{MIN}\end{cases}$
$\operatorname{H-MINIMAX}(\mathrm{s}, \mathrm{d})= \begin{cases}\operatorname{EVAL}(s) & \text { if } \operatorname{CUTOFF}-\operatorname{TEST}(s, d) \\ \max _{a} \mathrm{H}-\operatorname{MINIMAX}(\operatorname{RESULT}(s, a), \mathbf{d}+\mathbf{1}) & \text { if PLAYER}(s)=\text { MAX } \\ \min _{a} \operatorname{H}-\operatorname{MINIMAX}(\operatorname{RESULT}(s, a), \mathbf{d}+\mathbf{1}) & \text { if } \operatorname{PLAYER}(s)=\text { MIN }\end{cases}$

## Evaluation function

Estimates utility of game from truncated position

- Order terminal states in same manner
$\square$ Fast to compute
$\square$ For non-terminal states, correlated with the truth

Weighted linear combination of features
$\square$ independence assumption

$$
\operatorname{EVAL}(s)=w_{1} f_{1}(s)+\ldots+w_{n} f_{n}(s)=\sum_{i=1}^{n} w_{i} f_{i}(s)
$$

## Heuristic difficulties


(a) White to move

(b) White to move

## Cutoff tests - when to stop?

$\square$ At a fixed depth
Iterative deepening
$\square$ Report the result of the last IDS search that was fully completedHorizon effect
$\square$ Pushing off the inevitable
Quiescence search
$\square$ Stop at quiescent (quiet) positions
$\square$ Focus on non-quiescent positions

## Stochastic Games

## Stochastic games

$\square$ Stochastic games include an element of chance
$\square$ e.g. dice, unpredictable or random opponents
Example 1-player and 2-player stochastic games
$\square$ solitaire, minesweeper, backgammon, pacman

How do we find the optimal strategy in the presence of uncertainty?

## Stochastic games

How do we find the optimal strategy in this case?
$\square$ Change minimax tree to include chance nodes
probability of occurring


Compute average (expected) utility

- e.g. $1 / 2(20)+1 / 2(2)=11$


## Stochastic games

How do we find the optimal strategy in this case?
$\square$ Update minimax tree to include chance nodes


Compute average (expected) utility
-e.g. $1 / 2(20)+1 / 2(2)=11$

## Stochastic games

$\square$ 2-person stochastic game: MAX, MIN, chance nodes


## Stochastic games

For a 2-person stochastic game: MAX, MIN, chance


## Stochastic games

$\square$ For a 2-person stochastic game: MAX, MIN, chance


## Stochastic games

For a 2-person stochastic game: MAX, MIN, chance


## Example: Backgammon

$\square$ White rolls 6-5: (5-10, 5-11), (5-11,19-24), ...


## ExpectiMinimax

Update minimax strategy to compute weighted average (expected value) at chance nodesGives the expected value of a position/move
$\operatorname{EXPECTIMINIMAX}(\mathrm{s})= \begin{cases}\operatorname{UTILITY}(s) & \text { if TERMINAL-TEST}(s) \\ \max _{a} \operatorname{EXPECTIMINIMAX}(\operatorname{RESULT}(s, a)) & \text { if PLAYER }(s)=\text { MAX } \\ \min _{a} \operatorname{EXPECTIMINIMAX}(\operatorname{RESULT}(s, a)) & \text { if PLAYER }(s)=\text { MIN } \\ \sum_{r} P(r) \cdot \operatorname{EXPECTIMINIMAX}(\operatorname{RESULT}(s, r)) & \text { if PLAYER }(s)=\operatorname{CHANCE}\end{cases}$
$\mathrm{O}\left(b^{m} n^{m}\right)$ where n is number of distinct outcomes

## ExpectiMinimax in real-time

Recall, real-time games need to make fast decisions
For minimax, scale of the evaluation function doesn' $\dagger$ matter

For expectiMinimax, the scale is important


## Summary of adversarial games

Talked mostly about zero-sum games

- If I win then you lose

Minimax is optimal strategy but often too slow
$\square$ Alpha-beta pruning can increase max depth by factor of 2
$\square$ Implement evaluation function and cutoff-tests for early stop and guided search

Expectiminimax used for games with an element of chance/randomness

