## ADVERSARIAL SEARCH

## Today

$\square$ Reading

- AIMA Chapter 5.1-5.5, 5.7,5.8
$\square$ Goals
- Introduce adversarial games
- Minimax as an optimal strategy
$\square$ Alpha-beta pruning
$\square$ (Real-time decisions)


## Adversarial Games

$\square$ People like games!
$\square$ Games are fun, engaging, and hard-to-solve
$\square$ Games are amenable to study: precise, easy-torepresent state space


Game pieces found in a burial site in Southeast Turkey. Dated about 3000 BC

"Game of Twenty squares" discovered in a burial site in Ur. Dated about 2550-2400 BC


Backgammon is also among one of the oldest games still played today

## Adversarial Games

Two-player games have been a focus of AI as long as computers have been around

## Checkers



Solved: state space was completely mapped out!

## Backgammon and Chess



Computers can compete at a championship level mappal


Computers are still
 at an amateur club-level

## Adversarial Games

Humans and computers have different relative strengths in game play

```
    humans
```

good at evaluating the strength of a board
for a player

good at looking ahead in the game to find winning combinations of moves

## How humans play games

An experiment (by deGroot) was performed in which chess positions were shown to novice and expert players.
experts could reconstruct these perfectly novice players did far worse...

Random chess positions (not legal ones) were then shown to the two groups
experts and novices did just as
badly at reconstructing them!


## How computers play games



## Terminology

$\square$ deterministic vs. stochastic games
$\square$ initial state, successor function, goal test,...
$\square$ utility function: defines the final numeric value for a game that ends in terminal state $s$ for player $p$
$\square$ Chess: $+1,0,1 / 2$ for a win, loss, or draw
$\square$ zero-sum game: equal and opposite utilities

- If I win, you lose.

Chess: $0+1,1+0,1 / 2+1 / 2$
$\square$ policy: a function that maps from the set of states to the set of possible actions

## Branching factor and depth

On average, there are fewer than 40 possible moves that a chess player can make from any board configuration...


## Simplified representation for twoplayer games



## Minimax: an optimal strategy

$\square$ An optimal strategy is one that is at least as good as any other, no matter what the opponent does - If there's a way to force the win, it will $\square$ Will only lose if there's no other option
$\square$ Minimax is an optimal strategy assuming both players play optimally

## Minimax: an optimal strategy



What action should MAX take?

## Minimax: an optimal strategy


If I did this, then
he would do
that, but then I
would do that,
and then he
would do this...
$\operatorname{MINIMAX}(\mathrm{s})= \begin{cases}\operatorname{UTILITY}(s) & \text { if TERMINAL-TEST}(s) \\ \max _{a} \operatorname{MINIMAX}(\operatorname{RESULT}(s, a)) & \text { if PLAYER }(s)=\operatorname{MAX} \\ \min _{a} \operatorname{MINIMAX}(\operatorname{RESULT}(s, a)) & \text { if PLAYER }(s)=\operatorname{MIN}\end{cases}$

## Minimax: An Optimal Strategy

function Minimax-Decision(state) returns an action
$v \leftarrow$ Max-Value(state)
return the action in SUCCESSORS(state) with value $v$
function MAX-VALUE(state) returns a utility value
if Terminal-Test(state) then return Utility (state)
$v \leftarrow-\infty$
for $a, s$ in SUCCESSORS(state) do
$v \leftarrow \operatorname{Max}(v, \operatorname{Min}-\operatorname{Value}(s))$
return $v$
function Min-VALUE(state) returns a utility value
if Terminal-Test(state) then return Utility (state)
$v \leftarrow \infty$
for $a, s$ in SUCCESSORS(state) do
$v \leftarrow \operatorname{Min}(v, \operatorname{Max}-\operatorname{Value}(s))$
return $v$

## Minimax: Baby Nim

## 111/1

Take 1 or 2 at each turn
Goal: take the last match

## Minimax: Baby Nim



Take 1 or 2 at each turn Goal: take the last match





## Minimax



## Properties of Minimax

$\square$ Minimax performs depth-first exploration of game tree.
$\square$ Recall time complexity for DFS is $\mathrm{O}\left(\mathrm{b}^{m}\right)$

For chess, $b \approx 35, d \approx 100$ for "reasonable" games

- exact solution completely infeasible

How can we find the exact solution faster?


## Alpha-Beta Pruning

$\square$ Alpha-beta pruning: eliminate parts of game tree that don't affect the final result
$\square$ Consider a node $n$

- If a player has a better choice $m$ (at a parent or further up), then $n$ will never be reached
$\square$ Once we know enough about $n$ by looking at some successors we can prune it.


## Alpha-Beta Example

Do depth-first search until first leaf


## Alpha-Beta Example



Alpha-Beta Example


## Alpha-Beta Example



Alpha-Beta Example


## Alpha-Beta Example



Alpha-Beta Example


## Alpha-Beta Example

MAX


## Alpha-Beta pruning

function ALPHA-BETA-SEARCH(state) returns an action
inputs: state, current state in game
$v \leftarrow \operatorname{MAX}-\operatorname{VALUE}($ state $,-\infty,+\infty)$
return the action in SUCCESSORS(state) with value $v$
function MAX- $\operatorname{VALUE}($ state $, \alpha, \beta)$ returns a utility value
inputs: state, current state in game
$\alpha$, the value of the best alternative for MAX along the path to state
$\beta$, the value of the best alternative for MIN along the path to state
if Terminal-Test(state) then return Utility(state)
$v \leftarrow-\infty$
for $a, s$ in SUCCESSORS(state) do
$v \leftarrow \operatorname{MAX}(\mathrm{v}, \operatorname{MiN}-\operatorname{VALUE}(s, \alpha, \beta))$
if $v \geq \beta$ then return $v$
$\alpha \leftarrow \operatorname{MAX}(\alpha, \mathrm{v})$
return $v$
function MIN-VALUE(state, $\alpha, \beta$ ) returns a utility value
inputs: state, current state in game
$\alpha$, the value of the best alternative for MAX along the path to state
$\beta$, the value of the best alternative for MIN along the path to state
if Terminal-Test(state) then return Utility(state)
$v \leftarrow+\infty$
for $a, s$ in SUCCESSORS(state) do
$v \leftarrow \operatorname{Min}(\mathrm{v}, \operatorname{MAX}-\operatorname{VALUE}(s, \alpha, \beta))$
if $v \leq \alpha$ then return $v$
$\beta \leftarrow \operatorname{Min}(\beta, \mathrm{v})$
return $v$

## Why is it called alpha-beta?

$\square \alpha$ is the value of the best (i.e., highest-value) choice found so far at any choice point along the path for MAX

If $v$ is worse than $\alpha$, MAX will avoid it $\square$ prune that branch
 MIN

## Properties of $\alpha-\beta$

Pruning does not affect final resultHowever, effectiveness of pruning affected by order in which we examine successors

What do you do if you don't get to the bottom of the tree on time?

